Patterns for Session-Based Access Control

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ABSTRACT

The concept of session, the context under wh ich a user accesses resources is very important to apply access control. We present first the Contr olled Access Sess ion patter n for d escribing how sessions can lim it the ri ghts of a user . We then combine this pattern with t wo exi sting access control p atterns. First we consider a patter n for Session-Based Role-Based Access Control, intended for organizations in which job functions form the basis for privilege as signments. Then, we pres ent a Session-Based Attribute-Based Access Control pattern for organizations in which accesses are controlled based on valu es of user attributes and object properties. Since the general properties of thos e patterns have been described earlier we emphasiz e the additional effect of using sessions. The Controlled Access Session pattern can also be combined with other models of access control or used on its own.

Categories and Subject Descriptors

D.3.3 [**Programming Languages**]: L anguage Contructs and Features – *patterns*.

D.2.11 [**Software Engineering**]: S oftware Architectures – *patterns*.

General Terms

Security, Documentation, Algorithms, Management

Keywords

Access session, acce ss contr ol, attr ibute-based acce ss contr ol, session-based access control, security patterns

1. INTRODUCTION

It is important to develop s ystems where security has been considered at all s tages of de sign, which not only satis fy their functional specifications but also satisfy security requirements. To do this we need to start with high-level models that represent the security policies of the inst itution. There are three models currently used by most s ystems: the ac cess m atrix, t he Role-Based Access Control (RBAC) model, and the multilevel model.

One of the first security models was the access matrix. The basic

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access matrix [13] included the tuple $\{s, o, t\}$, where s indicates a subject or active entity, o is the protected object or resource, and t indicates the ty pe of access permitted. [Har76] proved security properties of this model u sing the s o-called HRU (Harrison-Ruzzo-Ullman) model. In that model users are allowed to delegate their rights (dis cretionary property, delegatable authorization), implying a tuple $\{s, o, t, f\}$, where f is a Boolean copy f lag indicating if the right is allowe d to be delegated or not. A predicate was added to the basic r ule to all ow content-based authorization [7], bec oming $\{s, o, t, p, f\}$, where p is the predicate (the predicate could also include environment variables). Patterns for the basic rule and for the tuple $\{s_{0}, t, p, f\}$ were given in [9][23]. The rule could also include the concept of Authorizer (a), becoming {a,s ,o,t,p,f} [8] (Explicitly G ranted A uthorization). RBAC [22] can be considered a special interpretation of the basic authorization model, where subjects are roles instead of individual users. We presented two varieties of RBAC pa tterns in [9] and [23]. Subsequently, several variations and extens ions of thes e models have appeared. We presented a variation called Metadata-Based Access Control, which later we renamed Attribute-Based Access Control (ABAC) [19][20].

ABAC can be seen in two ways:

- A s pecialization of the model {s _,o,t,p}, where p is a predicate which depends on attribute values.
- A variant where s and o are defined by descriptors which depend on attribute values.

In this paper we present a general pattern for a Controlled Access Session a s a building block and two p atterns combining this pattern with specific access contr ol models . The concept of session, the context under which a user accesses resources is very important to apply access control. We present first the Controlled Access Session pattern for describing how sessions can limit the rights of a user. We then combine this pattern with a pattern for Session-Based Role-Based Access Control, inte nded for organizations in which job functions form the basis for privilege assignments. Then, we present a Session-Based Attribute-Based Access Control pattern for organizations in which ac cesses a re controlled based on values of user attributes and object properties. Since the general properties of those patterns have been described earlier we emphasize the additional effect of us ing sessions. The Controlled Acces s Ses sion pattern can als o be combined with other models of ac cess control or used on its own. The pattern diagram of Fig ure 1 shows the relationship s between thes e patterns. For example, adding a condition to Basic Authorization results in Conte nt-Based Authorization, us ing the concept of session results in s ession-based models, and s o on. Note that RBAC is, in general, not delegatable. All these patterns define authorization r ules and they need a reference monitor f or their

enforcement; we don't show it in this diagram for simplicity (see [23] for the corresponding pattern). The double -lined patterns are the ones presented here. We assume the reader to know bas ic security concepts and these patterns are intended for system designers trying to add security to their designs.

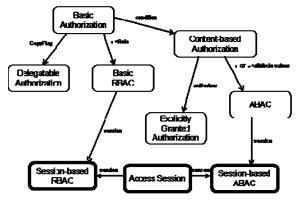


Figure 1. Relationships between access and control patterns

2. Controlled Access Session

Provide a context in which a subject (us er, s ystem) can access resources with different rights and without need to reauthenticate every time he accesses a new resource.

2.1 Example

Lisa is a secretary in a medical organization but sometimes she helps in the laboratory to perform patient tests. As a secretary she has access to patients ' information such as name, addre ss, SSN, etc. This is necess ary so she can bill them and their insurance companies. In the lab s he has access to anony mized patient test results. Combining the access es provided by her two jobs in one window she can ass ociate test results to names, which violates patient privacy.

2.2 Context

Any environment where we need to control access to computing resources and where us ers c an be clas sified ac cording to their jobs, groups, departments, assignments, or tasks.

2.3 Problem

A given user may be authorized to access a system because she needs to perform several functional ac tivities. However, for a particular access only those privileges should be active which are necessary to perform the intended task. This is an application of the principle of least-privilege and necessary to p revent the user from mis using the system (intention ally, accidentally by performing an error, or without knowledge and tricked to do so, for e xample through a Troj an Horse att ack). Additionally this would potentially restrict damage in case of session hijacking. A successfully attacking process would not have all privileges of a user available but only the active subset.

The following forces will affect the solution:

• Subjects may have many rights directly or indirectly through the execution contexts that they need for their tasks. Using all of them at one time may result in conflicts of interest and security violations. We need to restrict the use of those rights depending on the application or task the subject is performing.

- In the context of an interaction we can make the access to some functions implicit, thus facilitating the use of the system and preventing errors that may result in vulnerabili ties. For example, som e editor s or other tools could be implicitly available in some sessions.
- It is not convenient to make subjects reauthenticate every time they request a new r esource. Once the s ubject is authenticated, this condition s hould remain valid during the whole session.

2.4 Solution

Define a unit of interaction, a *session*, which has a limited lifetime, e.g. between login and logoff of a user or between the beginning and the end of a transaction. When a user logs on and after authentication, the session activates some execution contexts with only a subset of the authorizations she possesses. It should be the minimal s ubset which is needed for the user or transaction to perform the intended ta sk. Only thos e rights are ava ilable within the session. A subject can be in sever al sessions at the same time; however, in every s ession only the necessary rights are active.

Structure

Figure 2 shows the class model of the Access Session pattern. Classes **Subject** and **Session** have the obvious meaning. The class **ExecutionContext** contains the set of active rights that the user may use within the session.

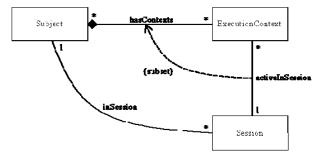


Figure 2. Class model for Access Session pattern

Figure 3 shows the use case Open (Activate) a session. A subject logs on and the logon interface authenticates it. The box with a double arrow indicates s ome authentication dialog or protoc ol. After the subject is authenticated, the interface creates a session object and returns a handle to the subject.

2.5 Implementation

Based on ins titution and application pol icies def ine w hich contexts (implying specific rights) should be us ed in e ach tas k and grant them to the corresponding subject. The rights should be selected using the least privilege principle and there should be no contexts with excessive rights, e.g. the administrator rights should be divided into smaller sets.

2.6 Example resolved

Lisa can log on a secretary or as a lab assistant but she cannot combine these activiti es in one ses sion. No w s he cannot r elate results to patient names.

2.7 Known uses

- Session Access is part of the RBAC standard proposal by NIST which later has been adopted by the American National S tandards Ins titute, I nternational Committee for Information T echnology Standards (ANSI/INCITS) as ANSI INCITS 359-2004 [10].
- Multics [Sum97] us ed exec ution context s (based on projects) to limit access right s. Ses sion Access is implemented in the s ecurity module CSAP [Dri03] of the Webocrat Sy stem in conju nction with an RBAC policy.
- Views in relational databases can be used to define sets of rights. Controlling the use of views by u sers can control their use of rights in sessions. This is done for example in Oracle and DB2, where SQL can be used to define restricted views [6].

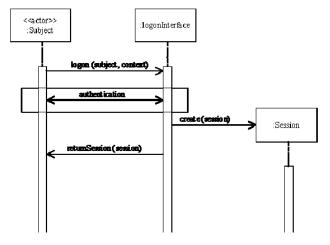


Figure 3. Sequence diagram for use case 'Open a session'

2.8 Consequences

This pattern has the following advantages:

- We can give to each context only the nee ded r ights according to its function and we can invoke in a ses sion only those contexts that are needed for a given task.
- We can exclude c ombinations of contexts that mig ht result in pos sible a ccess violati ons or conf licts of interest.
- Any functions can be made implicit in a session.
- Once a subje ct s tarts a ses sion it doesn't have to be reauthenticated. Its status is kept by the session.

Possible disadvantages:

- If we need to apply fi ne-grained access, it might be inefficient to include many contexts to perform complex activities.
- Using sessions may be confusing to the users.

2.9 Related patterns

The Access Session pattern is used in the Session-Based RBAC and ABAC patterns, discussed later.

The Session pattern of [26] created a session object that defined a namespace to hold all the variables that need to be referenced by many objects. P. Sommer lad remade this pattern as a Security Session [23], int ended to prevent a us er to b e r eauthenticated every time he ac cesses a new object. A pattern wit h a similar objective to the previous one is Abstract Session [21]: When an object's s ervices are invoked by clients, the s erver object may have to maintain state for each client. The server creates a session object that encapsulates s tate inf ormation for the client. The server returns a pointer to the ses sion object. However, none of these patterns, considers limitation of rights . Our p attern is an extension of those patterns, concentrating all its security functions and emphasizing the function of a session as a limiter of rights.

3. Session-Based Role-Based Access Control

Allow access to resources based on the role of the subject and limit the rights that can be applied at a given time based on the contexts (roles) defined by the access session.

3.1 Example

John is a developer in a project. He is also a project leader in another project. As a project leader he can evaluate the performance of the members of his project. He combines his two roles and adds several flattering evaluations about himself in the project where he is a developer. Later, his manager thinking that they came from the project leader of that project, gives John a big bonus.

3.2 Context

Any environment where we need to control access to computing resources, where users can be classified according to their jobs or their tas ks, and where we as sign rights to the roles needed to perform those tasks.

We assume the existence of a Session pattern that can be used for the solution.

3.3 Problem

In an organization a user may play several roles. However, for each access the user must act only within the authorizations of a single role (i.e. within the context of the role) or combinations of roles that do not violate ins titution policies. How do we force subjects to follow the policies of the institution when using their roles?

In addition to the forces defined for the Access Session pattern, the following forces apply to the solution:

- People in institutions have different needs for access to information, according to their functions. They may have several roles associated with specific functions or tasks.
- We want to help the institution to define precise access rights for its members so that the least privilege policy can be applied when they perform specific tasks.

• Users may have more than one role and we may want to enforce policies such as separation of duty, where a user cannot be in two or more specific roles in the same session.

3.4 Solution

A subject may have several roles. Each role collects the rights that a user can activate at a given moment (execution context), while a session controls the way of using roles and can enforce role exclusion at execution time.

Structure

The structure of the s ession-based RBAC is shown in the class diagram giv en in Figure 4. The class **Role** is an interm ediary between s ubject and object holding all authorizations a user possesses while playing the role and a cts here as an exe cution context. Within a **Session**, only a subset of the roles assigned to a **Subject** may be activated, i.e. only those necessary to perform the intended task. Roles may be composed according to a Composite pattern [11], where higher-level roles acquire (inherit) rights from the lower-level roles.

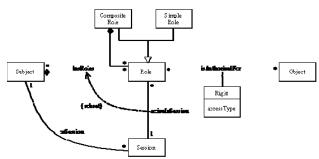


Figure 4. Class model for the Session-Based RBAC

Figure 5 shows a sequence diagram to request access to an object. A subject has a lready opened a session (See Figure 3) and he requests access to an object in a specific way (access type). The session uses the corre sponding Reference Monitor, which in turn checks if the rights of the session roles allow the access. If so, the access is permitted.

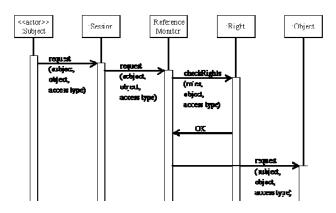


Figure 5. Sequence diagram to access an object

3.5 Implementation

See Section 5 for an example of a real implementation.

- Determine the roles the system should contain (role catalog), according to the user functions or tasks.
- Collect lis ts of incom patible r oles and use these lists when a session is s tarted (static cons traints). These constraints can be defined using OCL or some other formal language as additions to the class diagram of the pattern.
- Determine the number of roles which may be active within a session (dynamic constraints).
- When a user opens a s ession s he must dec lare what roles s he intends to use and the system will open the corresponding s ession or refus e to do so in case of conflicts.

3.6 Example resolved

When John logs on the project where he is a developer he only gets the rights for a developer and cannot add evaluations. When he logs on in the project where he is a project leader he can only evaluate the members of his group. He cannot combine his role rights in the same ses sion and now he only gets legitimate evaluations.

3.7 Known uses

The structure and dy namics of a s ession-based RBAC are implemented in the security module CSAP [5] of the Webocrat system. Webocrat is a portal supporting E-Democracy which was developed within the Europ ean Webocracy project (FP5-IST-1999-20364) between 2000-2003.

Views in relational databases can be used to define sets of rights. Controlling the use of views by roles can control the use of rights in sessions. In both Oracle and DB2 SQL can be used to define restricted views based on roles [6].

3.8 Consequences

In addition to the a dvantages mentioned for the Acc ess Session pattern, other advantages of this pattern are:

 Sessions may include all needed roles for those subjects authorized for some task.

- Users can activate more than one session at a time for functional flexibility (some tasks may require multiple roles).
- Fine-grained rights can be assigned to roles to enforce a need-to-know policy.
- When a s ession is open, we can exclude roles that violate institution policies.

Possible disadvantages include:

- Additional conceptual complexity to define which roles can be us ed together and which should be mutually exclusive.
- User conf usion if they have to use s everal r oles to perform their work.

3.9 Related Patterns

This pattern is a combinatio n of the Ses sion pattern described earlier a nd the RBAC pattern [23]. As indicated earlier , structuring of roles can be represented by a Composite pattern. A Reference Monitor pattern is needed to enforce the use of rights during execution.

4. Session-Based Attribute-Based Authorization

Allow access to resources based on the attributes of the subjects and the properties of the objects but limit the rights that c an be applied at a given time based on the context defined by the access session.

4.1 Example

Meili is a tee nager who likes movies and subscribes to several movie services through the Internet. She logs in a central portal where she can reach a variety of movies. S ometimes she gets movies that she find s offensive or inappropriate (pornographic, racist, plain stupid). She doesn't have much time to read details about the movies in advance and some of them don't even have good descriptions so reading about the movie s is not a good approach. She would like s ome kind of filter according t o h er characteristics and her prefer ences. Al so the portal may be breaking the law in making available to her some of these movies.

4.2 Context

Dynamic systems supporting a large set of objects and subjects in which the structure of the s ubjects changes rapidly, such as webbased information sy stems, e-government and e-business portals. In this e nvironment ther e is the need to control access to computing resources and the subjects may not be preregistered. We want to give access to resources based on characteristics of the subjects such as groups to which they belong, company for which they work, biological characteristics such as age or sex, or on characteristics of the objects, such as type of object, filtering rules, or payment requirements.

4.3 Problem

As indicated access may depend on the age or other attributes of a user. In this case, privilege assignments to the user cannot be done statically by a s ecurity adm inistrator but autom atically by the system based on the value of s ome of the attributes , e.g. "DateOfBirth" . As the user gets older or changes functions his authorization state changes automatically . Acc ess rights might even depend on an external attribute, such as "phy sical location" of a user in a mobile environment. In this case the authorization state changes automatically when the user moves around. At the object's side, metadata such as the s cope of a document, or the MPAA rating of a movie a re examples of proper ties. All these constraints can be applied through predicates in the rules [8], but it is difficult to have a variety of prepackaged rules for the typical cases.

The solution is constrained by the following forces:

- We need to limit the rights of s ubjects that are in a variety of groups or roles, or have special characteristics. Unrest ricted ac cess might allow policy or law violations.
- This control should not imply an extra burden for the security administrator or s ecurity vulnerabili ties may appear through administration errors.
- This control should not imply a significant performance overhead, or the system may not be practical to use.
- The environment is very dynamic and changes should be easy to make. Otherwise, the users will get annoyed and leave the system.

4.4 Solution

Access rights are based on the c omparison of values of s elected attributes of subjects and properties of objects (so called subject and object d escriptors). In this pattern descriptors are a construct to somehow "group" objects and su bjects dy namically, not explicitly by an administrator but implicitly by their attribute or property values. This grouping may result in unpredictable sets of rights that may violate security policies. A session delimits the rights that can be applied at a given moment; that is, the subject attributes define a context for access rights.

Structure

Figure 6 shows the class diagram for the s olution. A **Subject Descriptor** is for med by apply ing **Qualifiers** (>, +,...) to **Attribute Values** to define constraints such as 'age > 15'. A **Session** selects some specific attribute values as execution context that defines the **Subject** descriptor at this moment. S imilarly, objects are defined based on the values of selected attributes.

4.5 Implementation

See Section 5 for an example of a real implementation.

1) Select a n a ppropriate pac kage to convey the s ubject's credentials including attributes. Examples wou ld be attribute certificates [15][17] or Kerberos tickets.

2) Select an implementation to express the object's attributes. Candidates could be standards on meta-data resource discovery, such as the Dublin Core Metadata Initiative [DCM].

3) Define an enforce ment mec hanism for the rights defined in contexts. See for example [2].

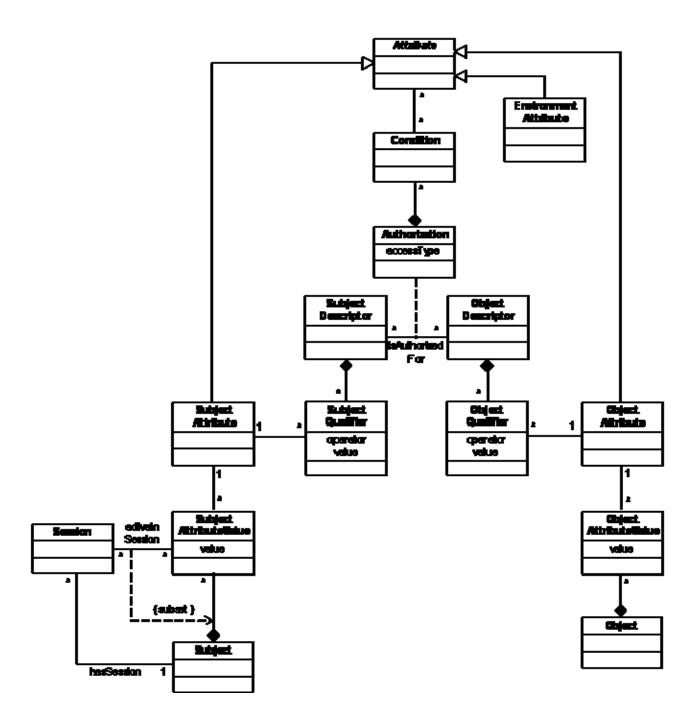


Figure 6. Class model for the Session-Based ABAC pattern

4.6 Example resolved

The portal implemented an ABAC model. Now when Meili opens a session she is given access to contexts with sets of preselected movies according to her preferences and restr icted a ccording to legal aspects and to the services she has paid for.

4.7 Known uses

Session-based ABAC is implemented as a n alternative to RBAC in the s ecurity m odule CSAP [5] of the W ebocrat system. A similar patter n is als o us ed in the author ization system of the .NET component framework [14] and in AAIs (authentication and authorization infrastructures), such as Permis [1] and Shibboleth [24].

The XM L s tandard XACML [4][16] uses attributes of s ubjects and objects for the specification of access control policies. As shown in the UCON_{ABC} [18], ABAC may also have potential for digital rights management.

4.8 Consequences

The advantages of this pattern include:

- The rights of subjects that belong to a variety of groups, roles, or have special attributes can be limited by restricting them to us e s pecific contexts s elected by sessions.
- This c ontrol do es not imply an extra burden for the security admini strator because the contexts can be defined by application n des igners according to application policies.
- This control does not imply a significant performance overhead because changing from one context to another just means changing a set of rights.
- Changes in access r estrictions can be easily accommodated by defining new contexts or deleting existing contexts.

Possible disadvantages are:

- Higher complexity. Although the contexts are defined by others, it is hard for administrators to know who has access to what.
- There might still be some performance overhead if we need to switch often between contexts.

4.9 Related Patterns

Figure 1 s hows the relations hip of t his patt ern to other access control patterns. As indicated credentials such as certificates are frequently used to request access [15].

5. Using session-based access control as a service

In this section we show by means of two sequence diagrams how the patterns described above can be embedded into a general authentication, author ization and access control s ervice. Such a service can be c alled by any application or p rocess having the need to authenticate the users and to provide session-based access control. In the following it is ass umed that the service provides both session-based RBAC and session-based ABAC and the client application requesting the service must chose between the two.

Figure 7 shows a sequence diagram for the interaction of a requesting client process and the s ession-based access control service. In order to hide the complexity of the subsystems, in the sequence diagram we use the Facade pattern [1 1] as a uniform interface for calling applications.

In order to be able to a ccess a resour ce, a valid session object must be requested by the calling application n (or us er proces s). This starts with some sort of initi alization process during which the client application first requests from the authentication facade of the security service an authentication service. In the example of Figure 7, a password service is returned but also other services may be available. Second is the request for an authoriz ation service. In the example, RBAC is returned, and the initialization phase is finished. N ext is the actual u ser authentication, role selection and the s ession es tablishment. During user authentication the client application provi des to the pas sword service
user-id, pw d>. The pas sword ser vice interacts with a userDM and in case of successful log-in a user object is created and a reference to the object (aUser) is returned to the calling client application.

A valid s ession can only be esta blished in the cas e the user application activates at least one role from the set of possible roles for the user. This starts by calling the method getAssignedRoles of the user object. In case of a valid userID all available roles for a particular us er are determ ined and returned by the role data module (RoleDM) and f or each role a transient role object is created by the RBAC service. Next from the set of possible roles the user s elects a sub set and the RBAC service calls the corresponding method to activate the roles.

At this stage the user object is authenticated and ha s a set of active r oles as signed. These ar e t he o nly prerequisites for establishing a ses sion. After receiving the r equest the ses sion service creates a valid session object for which the session-id is returned as a reference for the calling client process. Under a valid session-id the client may act under the context of the session by using the privileges of the selected roles.

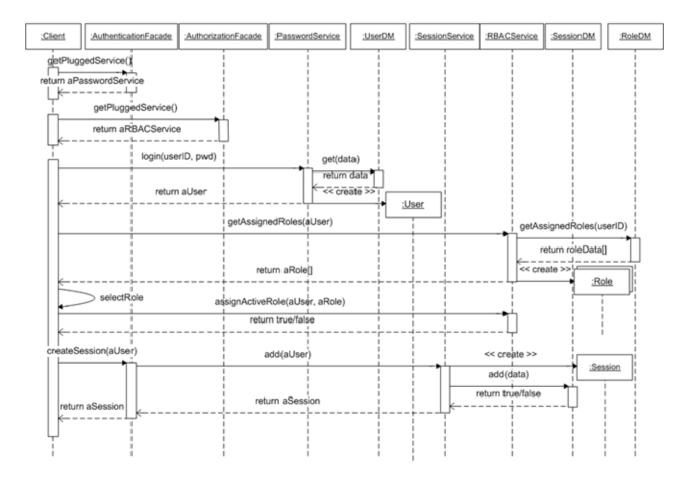


Figure 7. Session establishment

Figure 8 shows an attempt of a client process to access a resource within a valid session. The process starts with calling the method checkAccess with parameters session-id, object-id, operation, i.e. a request of a user wishing to access a certain object by using a predefined operation and this all wit hin the c ontext of an established session. F irst, the validity of the session is checked, then the session object is used by the RBACService in order to get the user' s active ro les within this session. N ext, the us er's permissions are d etermined by re trieving all the perm issions assigned to the active rol es. Fi nally, the RBACSer vice c hecks whether there is a permission for the tuple <object, operation>. In the case there is one, the access will be granted, otherwise denied.

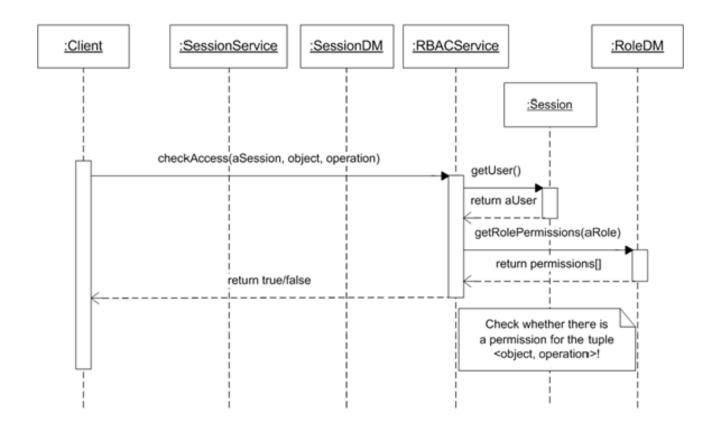


Figure 8. Permission approval

6. Conclusion

We have shown patterns to des cribe the effect of s essions on access contr ol models. We presented fir st the Access Session, which describes the basic concept of s ession as a limiter of rights. We then combined this pattern with the patterns of two access control models to show its effect on them. Finally we showed an example of a system using the last two patterns as a way to illustrate a real implementation.

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